

FD-350-96732600

SPECIFICATION

TO ALL WHOM IT MAY CONCERN:

BE IT KNOWN THAT WE, Hideo Miyake, a citizen of Japan residing at Kawasaki, Japan, Atsuhiko Suga, a citizen of Japan residing at Kawasaki, Japan and Yasuki Nakamura, a citizen of Japan residing at Kawasaki, Japan have invented certain new and useful improvements in

COMPUTER AND CONTROL METHOD OF THE COMPUTER

of which the following is a specification : -

COMPUTER AND CONTROL METHOD OF THE
COMPUTER

1. Field of the Invention

10 Generally, in a computer system, it is
necessary to deal with eventuality which is not
described in a program or which can not be described
in a program. In addition, it is necessary to
detect operations (abnormality, error, exception and
15 the like) other than intrinsic functions of hardware
and programs and to deal with the operations.
Further, it is necessary to provide communication
means between a hardware mechanism and system
software (especially OS), and it is necessary to
20 carry out scheduling conflicts of use requests for
shared hardware resources (processor, memory and the
like) so that these resources are used efficiently.

Therefore, from these, following interrupt
30 services are carried out as means for changing,
forcedly and dynamically, execution order of
instructions which are denoted as a program in the
computer system.

Generally, in the computer system, when
35 the interrupt occurs, control flow branches to the
interrupt service which is different from intrinsic
instruction execution order (which is denoted as the

5

10

15

25

35

following processes are necessary for realizing the precise interrupt.

These are (1) when executing an instruction which has a possibility of being interrupted, execution of following instructions are suppressed, (2), an interrupt which occurs first when executing a program is accepted and execution of following instructions and occurrence of interrupt are cancelled.

The operation of the above-mentioned case (1) can be realized by interlocking an instruction issuing part when the instruction which has a possibility of being interrupted is executed and the operation of the above-mentioned case (2) can be realized by providing an exception recovery mechanism such as checkpoint repair, history buffer, reorder buffer and future file and the like which are described in "Superscalar Microprocessor Design" by Mike Johnson (Prentice-Hall, inc., 1991).

In the following, an example of a conventional computer will be described. Fig.1 is a block diagram of a conventional computer. As shown in Fig.1, the conventional computer includes a memory 1, an instruction reading part 3, an instruction execution part 5, a register control part 6 and an interrupt control part 7. The instruction reading part 3 includes an instruction reading control part 9, an instruction word register (IR) 11 and a program counter 13. The instruction execution part 5 includes an instruction decode part 15, instruction issue checking parts 17, 18, OR circuits 19, 21, a selector 23, a load instruction execution part 25, a store instruction execution part 27, an arithmetical instruction execution part 29, an other instruction execution part 31, a floating-point load instruction execution part 33, a floating-point store instruction execution part 35.

00038406-000704

and a floating-point arithmetical instruction execution part 37.

5 The register control part 6 includes a register (PCSR) 39, a register (EPSR) 41, a register (PSR) 43, a general register 45 and a floating-point register 47. The interrupt control part 7 includes an interrupt control circuit 49.

10 The instruction reading part 3 reads an instruction word indicated by the program counter 13 and provides the instruction word to the instruction execution part 5. The instruction execution part 5 executes the provided instruction. The program counter 13 indicates an address of the instruction word.

15 The register control part 6 performs control for read and write of data by the instruction execution part 5, the instruction reading part 3 and the interrupt control part 7.

20 More precisely, the instruction reading control part 9 reads an instruction word indicated by the program counter 13 from the memory 1 and writes the word into the instruction word register 11. When a branch address is provided by the instruction execution part 5, or when a branch
25 address corresponding to interrupt or exception is provided from the interrupt control part 7, the branch address is written into the program counter 13. In the other cases, the program counter 13 is incremented so that next instruction word is
30 indicated. In addition, when the interrupt is detected while reading the instruction word, this is notified to the interrupt control circuit 49. When an instruction issue suppression notification is received from the instruction issue checking parts
35 17 and 18, the instruction reading operation is suspended until the instruction issue suppression notification is released.

FD-2230-50782600

The instruction decode part 15 analyses the instruction provided from the instruction register 11 and provides the instruction to a corresponding instruction execution part via the selector 23. When a load instruction is provided, it is provided to the load instruction execution part 25, when a store instruction is provided, it is provided to the store instruction execution part 27, when an arithmetical instruction is provided, it is provided to the arithmetical instruction execution part 29, and when a branch instruction is provided, it is provided to the other instruction execution part 31.

In addition, when a PCSR reading instruction, a PCSR writing instruction, an EPSR reading instruction or an EPSR writing instruction is provided, it is provided to the other instruction execution part 31. Further, when a floating-point load instruction is provided, it is provided to the floating-point load instruction execution part 33, when a floating point store instruction is provided, it is provided to the floating-point store instruction execution part 35, when a floating-point arithmetical instruction is provided, it is provided to the floating-point arithmetical instruction execution part 37.

When the selector 23 receives an instruction issue suppression notification from the instruction issue checking parts 17, 18, the selector suppresses providing an instruction to the instruction execution parts. Then, the instruction issue checking part 17 checks whether a focusing instruction is executable independently of an executing instruction. When the instruction can not be executed independently due to dependency, an instruction issue suppression notification is provided to the instruction reading control part 9

In the same way, the instruction issue checking part 18 checks whether an instruction which has a possibility of being interrupted is executed. When the instruction which has a possibility of being interrupted is executed, the instruction issue suppression notification is provided to the instruction reading control part 9 or the instruction decode part 15. When execution of the instruction which has a possibility of being interrupted ends, the instruction issue suppression notification is released.

25 When the store instruction is provided,
the store instruction execution part 27 obtains an
effective address from a value read from the general
register and writes data read from the general
30 register 45 to an area in the memory 1 corresponding
to the effective address. The data written into the
memory 1 is called "store data". When interrupt is
detected while executing the instruction, this is
notified to the interrupt control circuit 49. When
execution of the instruction ends, it is notified to
35 the instruction issue checking parts 17, 18.

When the arithmetical instruction is provided, the arithmetical instruction execution

part 27 performs an operation based on a value read from the general register 45 and writes the operation result to the general register 45. The data written into the memory 1 is called "store data". When interrupt is detected while executing the instruction, this is notified to the interrupt control circuit 49. When execution of the instruction ends, it is notified to the instruction issue checking parts 17, 18.

10 When the branch instruction is provided, the other instruction execution part 31 provides the branch address to the program counter 13. When the interrupt recovery instruction is provided, a value of the EPSR 41 is written into the PSR 43 and data read from the PCSR 39 is provided to the program counter 13 as the branch address.

15 When the PCSR reading instruction is provided, data held in the PCSR 39 is read and written into the general register 45, and data read from the general register 45 is written into the PCSR 39. When the PSR reading instruction is provided, data held in the PSR 43 is written into the general register 45. When the PSR writing instruction is provided, data read from the general register 45 is written into the PSR 43.

20 When the EPSR reading instruction is provided, data held in the EPSR 41 is read and written into the general register 45. When the EPSR writing instruction is provided, data read from the general register 45 is written into the EPSR 41. When interrupt is detected while executing the instruction, this is notified to the interrupt control circuit 49. When execution of the instruction ends, it is notified to the instruction issue checking parts 17, 18.

35 When the floating-point load instruction is provided, the floating-point load instruction

TOP SECRET

5

10

25

35

The interrupt control circuit 49 writes an instruction address indicating a return address from the interrupt into the PCSR 39, writes a state before interrupt into the EPSR 41, writes a state of the computer corresponding to the interrupt into the PSR 43 on the basis of interrupt notification provided from the instruction reading control part 9 or the instruction execution part 5. A branch address corresponding to the occurred interrupt is provided to the program counter 13.

In addition, when the exception recovery mechanism is provided for realizing the precise interrupt, although decrease of the effective rate of operation of hardware can be avoided, there is a problem in that circuit size increases.

It is an object of the present invention
30 to provide a computer and a control method of the
computer for processing an interrupt without
decreasing the effective rate of operation of
hardware.

a data holding part which holds data at a

According to the invention, when an interrupt is needed to be carried out, the data holding part holds data at the time. Therefore, it can be avoided to rerun an instruction which has been performed before occurrence of the interrupt so that operation after the interrupt can be continued efficiently. Thus, the interrupt can be processed without decreasing the effective rate of operation of hardware.

Other objects, features and advantages of the present invention will become more apparent from 15 the following detailed description when read in conjunction with the accompanying drawings, in which:

20 Fig.2 is a block diagram of a computer of
an embodiment 1 of the present invention;

25 Fig.4 is a flowchart showing an operation
performed by execution of an interrupt processing
program in the computer of the embodiment 1;

30 Fig.6 is table indicating bits of
effective flags and corresponding registers in ESR;

35 Figs.8 and 9 are flowcharts of a first
case showing an operation executed by an interrupt
processing program in the computer of the embodiment

Figs.10 and 11 are flowcharts of a second case showing an operation executed by an interrupt processing program in the computer of the embodiment

Fig.12 is a block diagram of the computer of an embodiment 3 of the present invention;

10

In the following, embodiments of the present invention will be described in detail with reference to figures. In the drawings, the same reference numerals are used to identify corresponding features.

[Embodiment 1]

20

25

35

instruction or an EDR writing instruction is provided, the instruction decode part 54 of the embodiment 1 provides each instruction to the other instruction execution part 52 via the selector 23.

5 The instruction issue checking part 51 checks whether an instruction which has a possibility of being interrupted is executed, and, when the instruction which has a possibility of being interrupted is executed and a focusing
10 instruction is an instruction which has a possibility of being interrupted, the instruction issue checking part 51 provides an instruction issue suppression notification to the instruction reading control part 9 or the instruction decode part 54.
15 When execution of the instruction which has a possibility of being interrupted ends, the instruction issue suppression notification is released.

 When an branch instruction is provided,
20 the other instruction execution part 52 provides a branch address to the program counter 13. When an interruption recovery instruction is provided, a value in the EPSR 41 is written into the PSR 43, and, data read from the PCSR 39 is provided to the
25 program counter 13 as a branch address.

 When an EPCR reading instruction is provided, data held in the EPCR 53 is read and written into the general register 45. When an EPCR writing instruction is provided, data read from the
30 general register 45 is written into the EPCR 53. When an ESR reading instruction is provided, data held in the ESR 55 is read and written into the general register 45. When an ESR writing instruction is provided, data read from the general
35 register 45 is written into the ESR 55.

 When an EAR reading instruction is provided, data held in the EAR 57 is read and

FD230-9648E660

written into the general register 45. When an EAR writing instruction is provided, data read from the general register 45 is written into the EAR 57.

When an EDR reading instruction is
5 provided, data held in the EDR 59 is read and written into the general register 45. When an EDR writing instruction is provided, data read from the general register 45 is written into the EDR 59. When interrupt is detected while executing the
10 instruction, this is notified to the interrupt control circuit 61. When execution of the instruction ends, it is notified to the instruction issue checking parts 17, 51.

The EPCR 53 is a register which holds an
15 instruction address of an instruction which causes an interrupt, and the ESR 55 is a register which holds data which indicates a factor of the interrupt wherein both of the registers are set when the interrupt occurs. The EAR 57 is a register which
20 holds an effective address of a load instruction, a store instruction, a floating-point load instruction, a floating-point store instruction in which the address is set when interrupt is occurred by a load instruction, a store instruction, a floating-point
25 load instruction, or a floating-point store instruction.

The EDR 59 is a register which holds store data used for executing a store instruction and a floating-point store instruction in which the store
30 data is set when interrupt is occurred by a store instruction or a floating-point store instruction.

The interrupt control part 61 writes an instruction address indicating a return address from interrupt into the PCSR 39, writes a state before
35 occurrence of the interrupt into the EPSR 41, writes a state of the computer corresponding to the interrupt into the PSR 43, writes an instruction

FD-230 (Rev. 5-22-64)

address of an instruction which causes the interrupt
(which will be called "interrupt cause instruction")
into the EPCR 53, and writes data indicating the
interrupt factor into the ESR 55 on the basis of
5 interrupt notification provided from the instruction
reading control part 9 or the instruction execution
part 50. When the interrupt cause instruction is
the load instruction, the store instruction, the
floating-point load instruction or the floating-
10 point store instruction, the effective address is
written into the EAR 57. When the interrupt cause
instruction is the store instruction or the
floating-point store instruction, the store data is
written into the EDR 59. In addition, a branch
15 address corresponding to the occurred interrupt is
provided to the program counter 13.

In the computer of the embodiment 1, in
the interrupt operation, the EPCR 53, the ESR 55,
the EAR 57 and the EDR 59 can be associated with
20 each instruction shown in Fig.3. The white circle
in Fig.3 means that the corresponding register holds
data when an interrupt occurs while an instruction
of the left side of the circle is executed. This is
also true for after-mentioned Fig.7.

25 That is, as shown in Fig.3, when an
interrupt occurs while an arithmetical instruction,
a floating-point arithmetical instruction or other
instruction is executed, an instruction address of
the interrupt cause instruction is held in the EPCR
30 53 as information necessary for recovering the
interrupt (or restarting), and data indicating the
interrupt factor is held in the ESR 55. When an
interrupt occurs while an load instruction, a
floating-point load instruction is executed, further,
35 an effective address is held in the EAR 57 as
information necessary for recovering the interrupt
(or restarting). In addition, when an interrupt

TO 2220 9643E660

occurs while a store instruction or a floating-point store instruction is executed, further, store data is held in the EDR 59 as information necessary for recovering the interrupt (or restarting).

5 In the computer of the embodiment 1, recover from an interrupt (or restart) can be easily realized by reading information held as shown in Fig.3 when an interrupt is processed.

10 In the following, an operation performed by execution of an interrupt processing program in the computer of the embodiment 1 will be described with reference to a flowchart shown in Fig.4. In step S1, context is saved. In step S2, the other instruction execution part 52 reads interrupt factor
15 data from the ESR 55 according to a provided ESR reading instruction. Next, in step S3, the other instruction execution part 52 refers to the interrupt factor data and determines whether the interrupt is caused by a load instruction or a
20 floating-point load instruction. When the other instruction execution part 52 determines that the interrupt is caused by a load instruction or a floating-point load instruction, the process goes to step S4, and, when the other instruction execution
25 part 52 does not determine that the interrupt is caused by a load instruction or a floating-point load instruction, the process goes to step S10.

 The other instruction execution part 52 reads an instruction address from the EPCR 53
30 according to an EPCR reading instruction in step S4, and reads an effective address from the EAR 57 according to an EAR reading instruction in step S5. In step S6, an interrupt service is performed based on the above-mentioned instruction address, the
35 interrupt cause data and the effective address in step S6.

Next, when the interrupt service ends, the

TOP SECRET

context is recovered in step S7, and the computer returns from interrupt processing program operation according to an interrupt return instruction in step S8.

5 In step S10, the other instruction execution part 52 refers to the interrupt factor data and determines whether the interrupt is caused by a store instruction or a floating-point store instruction. When the other instruction execution
10 part 52 determines that the interrupt is caused by a store instruction or a floating-point store instruction, the process goes to step S11, and, when the other instruction execution part 52 does not
15 determine that the interrupt is caused by a store instruction or a floating-point store instruction, the process goes to step S20.

 The other instruction execution part 52 reads an instruction address from the EPCR 53 according to EPCR reading instruction in step S11,
20 and reads an effective address from the EAR 57 according to an EAR reading instruction in step S12, and reads store data from the EDR 59 according to an EDR reading instruction in step S13. Then, in step
25 S14, the interrupt service is performed based on the above-mentioned instruction address, the interrupt cause data, the effective address and the store data, and the process goes to step S7.

 The other instruction execution part 52 reads an instruction address from the EPCR 53
30 according to an EPCR reading instruction in step S20, and goes to step S14.

 As mentioned above, according to the computer of the embodiment 1, since the computer includes the EPCR 53, the ESR 55, the EAR 57 and the
35 EDR 59 where necessary data for returning to original process from the interrupt processing operation is held in these registers when interrupt

FIG. 10: S10-S14

occurs, and since the original process continues according to the held data after the interrupt processing ends, an interrupt can be performed with a simple configuration without decreasing effective
5 rate of operation of hardware.

[Embodiment 2]

Fig.5 is a block diagram of a computer of the embodiment 2 of the present invention. As shown in Fig.5, the computer of the embodiment 2 of the
10 present invention is almost the same as the computer of the embodiment 1 shown in Fig.2 except that an effective flag (ESFR) 71 is provided in a register control part 69, and an EPCR 73, an ESR 75, an EAR 77 and an EDR 79 each of which includes a plurality
15 of registers are included.

Each of the effective flag (ESFR) 71, EPCR 73, ESR 75, EAR 77 and EDR 79 is connected to an other instruction execution part 67 included in an instruction execution part 63 and to an instruction
20 control circuit 81 included in an interrupt control part 80. In addition, different from the embodiment 1, only one instruction issue checking part 65 is provided and the OR circuits 19, 20 are not included in the instruction execution part 63.

25 In the computer of the embodiment 2 which has the above-mentioned configuration, an instruction decode part 16 analyzes an instruction provided from the instruction word register 11 and provides the instruction to an corresponding
30 instruction execution part via the selector 23. For example, when a branch instruction, an interrupt recovery instruction, an ESFR reading instruction or an ESFR writing instruction is provided, each instruction is provided to the other instruction
35 execution part 67.

When a branch instruction is provided, the other instruction execution part 67 provides a

09938496-082704

branch address to the program counter 13. When an interrupt recovery instruction is provided, a value in the EPSR 41 is written into the PSR 43 and data read from the PCSR 39 is provided to the program counter 13 as the branch address.

When an ESFR reading instruction is provided, data held in the ESFR 71 is read and written into the general register 45. When an ESFR writing instruction is provided, data read from the general register 45 is written into the ESFR 71. When an interrupt is detected while executing the instruction, this is notified to the interrupt control circuit 81. When execution of the instruction ends, it is notified to the instruction issue checking part 65.

The effective flag (ESFR) 71 is an effective flag of the ESR 75 which holds interrupt factor data where each bit of the flag corresponds to each register and the flag indicates whether the interrupt factor data is held in the register. The effective flag 71 is set when an interrupt occurs.

When the ESR 75 includes seven registers in parallel, each bit from 0th bit to 6th bit corresponds to each register as shown in Fig.6.

As shown in Fig.6, for example, the 0th bit corresponds to the 0th ESR register in which the 0th bit becomes 1 when the interrupt factor data is held in the 0th ESR register.

The interrupt control part 81 writes an instruction address indicating a return address from interrupt into the PCSR 39, writes a state before interrupt into the EPSR 41, writes a state of the computer corresponding to the interrupt into the PSR 43, writes an instruction address of the interrupt cause instruction into the EPCR 73, writes data indicating interrupt factor to ESR 75, and writes holding state of the data indicating interrupt

00000000-00000000

factor into ESFR 71 on the basis of interrupt notification provided from the instruction reading control part 9 or the instruction execution part 63.

When the interrupt cause instruction is the load instruction, the store instruction, the floating-point load instruction or the floating-point store instruction, the effective address is written into the EAR 77. When the interrupt cause instruction is the store instruction or the floating-point store instruction, the store data is written into the EDR 79. In addition, a branch address corresponding to the occurred interrupt is provided to the program counter 13.

In the computer of the embodiment 2, for example, when the EPCR 73 includes parallel seven registers from EPCR0 to EPCR6, the ESR 75 also includes parallel seven registers from ESR0 to ESR6, the EAR 77 includes parallel four registers from EAR3 to EAR6, and the EDR 79 includes parallel two registers of EDR5 and EDR6, these registers can be associated with each instruction as shown in Fig.7.

Next, operation executed by an interrupt processing program in the computer of the embodiment 2 will be described with reference to Figs.8-11 for two cases in which the ESFR 71 is used and the ESFR 71 is not used. First, operation in the case where the ESFR 71 is not used will be described with reference to Figs.8 and 9.

As shown in Fig.8, contest is saved in step S1. Then, the other instruction execution part 67 read interrupt factor data 0 from the ESR 75 according to a provided ESR reading instruction in step S2. In step S3, the other instruction execution part 67 judges whether the interrupt factor data 0 is held. When the other instruction execution part 67 judges that the interrupt factor data 0 is held, the process goes to step S100, and

FIG. 8

when the other instruction execution part 67 does not judge that the interrupt factor data 0 is held, the process goes to step S4.

5 In step S100, the other instruction execution part 67 reads an instruction address 0 from the EPCR 73 according to the EPCR reading instruction. Then, the interrupt service is performed based on the interrupt factor data 0 and the instruction address 0, and the process goes to
10 step S4.

In step S4, the other instruction execution part 67 read interrupt factor data 1 from the ESR 75 according to a provided ESR reading instruction. In step S5, the other instruction
15 execution part 67 judges whether the interrupt factor data 1 is held. When the other instruction execution part 67 judges that the interrupt factor data 1 is held, the process goes to step S110, and when the other instruction execution part 67 does
20 not judge that the interrupt factor data 1 is held, the process goes to step S6.

In step S110, the other instruction execution part 67 reads an instruction address 1 from the EPCR 73 according to the EPCR reading
25 instruction. Then, the interrupt service is performed based on the interrupt factor data 1 and the instruction address 1 , and the process goes to step S6.

In step S6, the other instruction
30 execution part 67 reads interrupt factor data 2 from the ESR 75 according to a provided ESR reading instruction. In step S7, the other instruction execution part 67 judges whether the interrupt factor data 2 is held. When the other instruction
35 execution part 67 judges that the interrupt factor data 2 is held, the process goes to step S120, and when the other instruction execution part 67 does

TOC230 9648650

In step S120, the other instruction execution part 67 reads an instruction address 2 from the EPCR 73 according to the EPCR reading instruction. Then, the interrupt service is performed based on the interrupt factor data 2 and the instruction address 2, and the process goes to step S8.

In step S8, the other instruction execution part 67 reads interrupt factor data 3 from the ESR 75 according to a provided ESR reading instruction. In step S9, the other instruction execution part 67 judges whether the interrupt factor data 3 is held. When the other instruction execution part 67 judges that the interrupt factor data 3 is held, the process goes to step S130, and when the other instruction execution part 67 does not judge that the interrupt factor data 3 is held, the process goes to step S10.

In step S130, the other instruction execution part 67 reads an instruction address 3 from the EPCR 73 according to the EPCR reading instruction, and the other instruction execution part 67 reads an effective address 3 from the EAR 77 according to the EAR reading instruction. Then, the interrupt service is performed based on the interrupt factor data 3, the instruction address 3 and the effective address 3, and the process goes to step S10.

In step S10, the other instruction execution part 67 reads interrupt factor data 4 from the ESR 75 according to a provided ESR reading instruction. In step S11, the other instruction execution part 67 judges whether the interrupt factor data 4 is held. When the other instruction execution part 67 judges that the interrupt factor

5 In step S140, the other instruction
execution part 67 reads an instruction address 4
from the EPCR 73 according to the EPCR reading
instruction, and the other instruction execution
part 67 reads an effective address 4 from the EAR 77
10 according to the EAR reading instruction. Then, the
interrupt service is performed based on the
interrupt factor data 4, the instruction address 4
and the effective address 4 in step S142, and goes
to step S12.

15 In step S12, the other instruction
execution part 67 reads interrupt factor data 5 from
the ESR 75 according to a provided ESR reading
instruction. In step S13, the other instruction
execution part 67 judges whether the interrupt
20 factor data 5 is held. When the other instruction
execution part 67 judges that the interrupt factor
data 5 is held, the process goes to step S150, and
when the other instruction execution part 67 does
not judge that the interrupt factor data 5 is held,
25 the process goes to step S14 in Fig.9.

In step S150, the other instruction
execution part 67 reads an instruction address 5
from the EPCR 73 according to the EPCR reading
instruction, and the other instruction execution
30 part 67 reads an effective address 5 from the EAR 77
according to the EAR reading instruction. In step
S152, the other instruction execution part 67 reads
store data 5 from the EDR 79 according to the EDR
reading instruction. Then, the interrupt service is
35 performed based on the interrupt factor data 5, the
instruction address 5, the effective address 5 and
the store data 5 in step S153, and goes to step S14.

5
10

15
20

25

30

35

5

10

20

25

35

5 process goes to step S130.

10 execution part 67 reads interrupt factor data 2 from
the ESR 75 according to ESR reading instruction.

15 In step S6, the other instruction
execution part 67 determines whether the third bit
of the effective flag 71 is 0. When the other
instruction execution part 67 determines that the
third bit of the effective flag 71 is 0, the process
20 goes to step S7 in Fig.11, and when the other
instruction execution part 67 determines that the
third bit of the effective flag 71 is not 0, the
process goes to step S140.

25 In step S140, the other instruction
execution part 67 reads an instruction address 3
from the EPCR 73 according to EPCR reading
instruction. In step S141, the other instruction
execution part 67 reads interrupt factor data 3 from
the ESR 75 according to ESR reading instruction. In
30 step S142, the other instruction execution part 67
reads an effective address 3 from the EAR 77
according to EAR reading instruction. Then,
interrupt service is performed based on the
instruction address 3, the interrupt factor data 3
35 and the effective address 3 in step S143, and the
process goes to step S7.

In step S7, the other instruction

execution part 67 determines whether the 4th bit of the effective flag 71 is 0. When the other instruction execution part 67 determines that the 4th bit of the effective flag 71 is 0, the process
5 goes to step S8, and when the other instruction execution part 67 determines that the 4th bit of the effective flag 71 is not 0, the process goes to step S150.

In step S150, the other instruction
10 execution part 67 reads an instruction address 4 from the EPCR 73 according to EPCR reading instruction. In step S151, the other instruction execution part 67 reads interrupt factor data 4 from the ESR 75 according to ESR reading instruction. In
15 step S152, the other instruction execution part 67 reads an effective address 4 from the EAR 77 according to EAR reading instruction. Then, interrupt service is performed based on the instruction address 4, the interrupt factor data 4
20 and the effective address 4 in step S153, and the process goes to step S8.

In step S8, the other instruction execution part 67 determines whether the 5th bit of the effective flag 71 is 0. When the other
25 instruction execution part 67 determines that the 5th bit of the effective flag 71 is 0, the process goes to step S9, and when the other instruction execution part 67 determines that the 5th bit of the effective flag 71 is not 0, the process goes to step
30 S160.

In step S160, the other instruction execution part 67 reads an instruction address 5 from the EPCR 73 according to EPCR reading instruction. In step S161, the other instruction
35 execution part 67 reads interrupt factor data 5 from the ESR 75 according to ESR reading instruction. In step 162, the other instruction execution part 67

TOP SECRET

5

10

25

35

As mentioned above, according to the computer of the embodiment 2, the computer includes the EPCR 73, the ESR 75, the EAR 77 and the EDR 79 each of which includes a plurality of registers in parallel, where necessary data for returning to original process from interrupt processing operation is held in these registers in parallel when interrupt occurs. Therefore, an instruction next to an instruction which has a possibility of being interrupted can be issued without waiting for completion of execution of the instruction which has a possibility of being interrupted. Therefore, effective rate of operation of hardware can be further improved.

[Embodiment 3]

Fig.12 is a block diagram of the computer of the embodiment 3 of the present invention. As shown in Fig.12, the configuration of the computer of the embodiment 3 is almost the same as that of the embodiment 1 except that an EAR 95 included in a register control part 93 is connected to a load instruction execution part 85, a store instruction execution part 87, a floating-point load instruction execution part 89 and a floating-point store instruction execution part 91 which are included in an instruction execution part 83, and an EDR 97 is connected to a store instruction execution part 87 and a floating-point store instruction execution part 91, and an interrupt control circuit 99 included in an interrupt control part 100 is connected to the EPCR 53 and the ESR 55 and is not connected to the EAR 95 and the EDR 97.

In the computer of the embodiment 3 having the above-mentioned configuration, when a load instruction is provided, the load instruction execution part 85 obtains an effective address from a value read from a general register 45, reads data

5

10

25

35

writes data read from the floating-point register 47 into an area in the memory 1 corresponding to the effective address. When interrupt is detected while executing the instruction, this is notified to the interrupt control circuit 49, and the effective address is written into the EAR 95 and the store data is written into the EDR 97. When execution of the instruction ends, it is notified to the instruction issue checking parts 17, 51.

As for the other operation, the computer of the embodiment 3 operates according to the flowchart shown in Fig.4 like the computer of the embodiment 1.

In the same way as the computer of the embodiment 1, according to the computer of the embodiment 3 of the present invention, an interrupt processing can be performed with a simple configuration without decreasing effective rate of operation of hardware.

[Embodiment 4]

Fig.13 is a block diagram of the computer of the embodiment 4 of the present invention. As shown in Fig.13, the configuration of the computer of the embodiment 4 is almost the same as that of the embodiment 2 shown in Fig.5 except that an EAR 105 included in a register control part 103 is connected to a load instruction execution part 86, a store instruction execution part 88, a floating-point load instruction execution part 90 and a floating-point store instruction execution part 92 which are included in an instruction execution part 101, and an EDR 107 is connected to a store instruction execution part 88 and a floating-point store instruction execution part 92, and an interrupt control circuit 109 included in an interrupt control part 110 is connected to the ESFR 71, the EPCR 73 and the ESR 75 and is not connected

In the computer of the embodiment 4 having the above-mentioned configuration, when a load instruction is provided, the load instruction

15 When the store instruction is provided,
the store instruction execution part 88 obtains an
effective address from a value read from the general
register 45 and writes data read from the general
register 45 into an area in the memory 1
20 corresponding to the effective address. When
interrupt is detected while executing the
instruction, this is notified to the interrupt
control circuit 109, and the effective address is
written into the EAR 105 and the store data is
25 written into the EDR 107. When execution of the
instruction ends, it is notified to the instruction
issue checking part 65.

When the floating-point load instruction is provided, the floating-point load instruction execution part 90 obtains an effective address from a value read from the general register 45, reads data from an area in the memory 1 corresponding to the effective address, and writes the result into the floating-point register 47. When interrupt is detected while executing the instruction, this is notified to the interrupt control circuit 109, and the effective address is written into the EAR 105.

When the floating-point store instruction is provided, the floating-point store instruction execution part 92 obtains an effective address from a value read from the general register 45, and writes data read from the floating-point register 47 into an area in the memory 1 corresponding to the effective address. When interrupt is detected while executing the instruction, this is notified to the interrupt control circuit 109, and the effective address is written into the EAR 105 and the store data is written into the EDR 107. When execution of the instruction ends, it is notified to the instruction issue checking part 65.

20 In the same way as the computer of the
embodiment 2, according to the computer of the
embodiment 4 of the present invention, the interrupt
processing can be performed with a simple
configuration without decreasing effective rate of
25 operation of hardware.

As mentioned above, according to the computer and the control method of the present invention, when interrupt becomes necessary, data at the time is held so that operation after the
30 interrupt can be continued effectively by avoiding rerunning an instruction before the interrupt. Thus, the interrupt processing can be performed without decreasing the rate of operation of hardware.

In addition, according to the present
35 invention, a plurality of data holding parts may be
provided in parallel so that a plurality of
instructions which may occur interruptions can be

executed in parallel. Therefore, instructions can
be executed further effectively. In addition, by
providing the effective flags which indicates
whether the data holding parts hold data, the
5 operation can be performed more speedily.

The present invention is not limited to
the specifically disclosed embodiments, and
variations and modifications may be made without
departing from the scope of the invention.

10

15

20

25

30

35

09928496-082704